### Introducing 3-Path Domination

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# Graphs - Definitions

#### Graph

G = (V, E) is a graph G with a set of vertices V(G) and set of edges E(G). We will work with *simple graphs* containing no loops or multiple edges.

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### Domination

#### Dominating Set

The set  $S \subseteq V(G)$  is a dominating set of G if for every vertex  $v \in V(G)$ , either  $v \in S$  or v is adjacent to a vertex in S. Every  $v \in S$  can be thought of as a guard that dominates adjacent rooms (vertices are rooms, edges are hallways).

#### Domination Number

The minimum cardinality of a dominating set, denoted  $\gamma(G)$ .

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### Domination

### Minimum Dominating Set Example

A case where  $\gamma(G) = 2$ . Every guard can watch every adjacent room, including theirs.



### Paired-Domination Introduced by Haynes and Slater in 1998

### Paired-Dominating Set

A dominating set where the induced subgraph on the set contains a perfect matching. In the sense of the guard analogy, every guard has another guard watching their back.

#### Paired-Domination Number

The minimum cardinality of a paired-dominating set, denoted  $\gamma_{pr}(G)$ , is always even.

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# 3-Path Domination

#### 3-Path Dominating Set

Let  $Q_i$  represent a path on 3 vertices. Then we define a 3-path dominating set of G to be  $S = \{Q_1, Q_2, \ldots, Q_k\}$  such that the vertex set  $V(S) = V(Q_1) \cup V(Q_2) \cup \cdots \cup V(Q_k)$  is a dominating set.

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# 3-Path Domination Number

#### 3-Path Domination Number

The minimum number of 3-paths needed to dominate a graph G, denoted  $\gamma_{P_3}(G)$ .

#### Justification

A vertex in a 3-path dominating set can be in more than one 3-path. So, the number of vertices in the set does not directly translate to the number of 3-paths.

# 3-Path Domination Number

#### Example

Below is an example of minimum 3-path dominating set. Although 9 vertices are being used, we have 4 3-paths.



# What is minimal?

#### Minimal

A dominating set S is said to be minimal if for a vertex  $v \in S$ ,  $S \setminus v$  is not a dominating set. A minimal dominating set is not a proper subset of any other dominating set.

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# The 3-Path Domination Problem is NP-complete

#### Theorem

Deciding for a given graph H and positive integer K such that  $3K \leq |V(H)|$ , "Is  $\gamma_{P_3}(H) \leq K$ ?" is NP-complete.

- In Haynes and Slater's paper, Paired Domination in Graphs, we get the result that paired domination is NP-complete.
- We generalized this result to show that the 3-path domination problem is also NP-complete.

# Upper Bound

#### Lemma

For a connected graph G, there exists an edge-disjoint minimal 3-path dominating set.



Figure: Two 3-paths,  $\{v_2, v_3, v_4\}$  and  $\{v_3, v_4, v_5\}$  dominating the path graph  $P_6$ .

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### Upper Bound

By the above lemma, we have  $\gamma_{P_3}(G) \leq \left| \frac{|E(G)|}{2} \right|$ .

### Trees

#### Tree

A tree  $T_n$  is an acyclic, connected graph on n vertices. Every tree has n - 1 edges.

### Spanning Tree

A spanning tree  $T_G$  of a graph G is a tree with vertex set V(G)and edge set  $E(T_G) \subseteq E(G)$ .



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A graph G with n = 30 vertices.



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# A graph G with n = 30 vertices.



# The Spanning Tree $T_G$ of G.



# A Small Example



Figure: An example where the blue path is no longer needed when an edge is added.

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# A More Explicit Bound

#### Theorem

For a connected graph G on n vertices,  $\gamma_{P_3}(G) \leq \left|\frac{n-1}{2}\right|$ .

#### Proof.

Let G be a connected graph on n vertices and  $T_G$  be its spanning tree. We know  $\gamma_{P_3}(T_G) \leq \lfloor \frac{n-1}{2} \rfloor$ . In order to obtain G from  $T_G$ , we need to add our missing edges, namely  $E(G) \setminus E(T_G)$ . Observe, adding edges to a graph will either keep the domination number the same or cause it to decrease. So,

$$\gamma_{P_3}(G) \le \gamma_{P_3}(T_G) \le \left\lfloor \frac{n-1}{2} \right\rfloor$$

# 3-Path Domination vs. Paired Domination

#### Theorem

For any connected graph G on  $n \geq 3$  vertices,  $\gamma_{P_3}(G) \leq \frac{\gamma_{Pr}(G)}{2}$ .

#### Sketch of Proof.

Let *D* be a minimum paired-dominating set of *G*. Every pair of vertices in *D* has a neighbor that is not in *D*. So we can count every pair and a neighbor as a 3-path. So,  $\gamma_{P_3}(G) \leq \frac{\gamma_{PT}(G)}{2}$ .

# 3-Path Domination vs. Paired Domination $_{\rm Example}$



Figure: The  $\gamma_{pr}$ -set from above turned into a  $\gamma_{P_3}$ -set

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## Conjecture

#### Theorem (Haynes et al.)

For any connected graph G on  $n \ge 6$  vertices and  $\delta(G) \ge 2$ ,  $\gamma_{pr}(G) \le \frac{2n}{3}$ .

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# Additional Results

- Comparisons between  $\gamma(G)$ ,  $\gamma_{pr}(G)$ , and  $\gamma_{P_3}(G)$
- Tighter bounds for certain families of graphs, such as:
  - Caterpillar Trees
  - Graphs with a Hamiltonian Path
- Closed formulas for certain families of graphs, such as:
  - Paths or Cycle Graphs
  - Harary graphs  $H_{k,n}$  where k is even
  - Banana Trees
- Algorithms for finding a minimal 3-path dominating set in trees

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